Module outline

1 Introduction to databases
2 Security requirements
3 Data disclosure and inference
4 Multilevel security databases
5 Data mining
(Relational) Databases

- Structured, queryable collection of data (records)
- Each record consists of fields (elements)
- Structure (schema) set by database administrator
- Database management system (DBMS) provides support for queries and management
- Most popular DBMS is based on relational model
- Stores records in one or multiple tables (relations)
  - Table has named columns (attributes) and rows (tuples)
  - Individual tables can have relationships between them

Records

<table>
<thead>
<tr>
<th>Name</th>
<th>First</th>
<th>Address</th>
<th>City</th>
<th>State</th>
<th>Zip</th>
<th>Airport</th>
</tr>
</thead>
<tbody>
<tr>
<td>ADAMS</td>
<td>Charles</td>
<td>212 Market St.</td>
<td>Columbus</td>
<td>OH</td>
<td>43210</td>
<td>CMH</td>
</tr>
<tr>
<td>ADAMS</td>
<td>Edward</td>
<td>212 Market St.</td>
<td>Columbus</td>
<td>OH</td>
<td>43210</td>
<td>CMH</td>
</tr>
<tr>
<td>BENCHLY</td>
<td>Zeke</td>
<td>501 Union St.</td>
<td>Chicago</td>
<td>IL</td>
<td>60603</td>
<td>ORD</td>
</tr>
<tr>
<td>CARTER</td>
<td>Marlene</td>
<td>411 Elm St.</td>
<td>Columbus</td>
<td>OH</td>
<td>43210</td>
<td>CMH</td>
</tr>
<tr>
<td>CARTER</td>
<td>Beth</td>
<td>411 Elm St.</td>
<td>Columbus</td>
<td>OH</td>
<td>43210</td>
<td>CMH</td>
</tr>
<tr>
<td>CARTER</td>
<td>Ben</td>
<td>411 Elm St.</td>
<td>Columbus</td>
<td>OH</td>
<td>43210</td>
<td>CMH</td>
</tr>
<tr>
<td>CARTER</td>
<td>Lisabeth</td>
<td>411 Elm St.</td>
<td>Columbus</td>
<td>OH</td>
<td>43210</td>
<td>CMH</td>
</tr>
<tr>
<td>CARTER</td>
<td>Mary</td>
<td>411 Elm St.</td>
<td>Columbus</td>
<td>OH</td>
<td>43210</td>
<td>CMH</td>
</tr>
</tbody>
</table>

Relations

NAME-ZIP

ZIP-AIRPORT

(attributes are missing)
Database queries

- Most popular query language is SQL
  - SELECT Address FROM NAME-ZIP
    WHERE (Zip = '43210') AND (Name = 'ADAMS')
    - Prints address of family in relation NAME-ZIP whose zip code is 43210 and whose name is Adams
  - SELECT Name, Airport
    FROM NAME-ZIP, ZIP-AIRPORT
    WHERE NAME-ZIP.Zip = ZIP-AIRPORT.Zip
    - Prints each person’s last name and his/her airport by joining relations NAME-ZIP and ZIP-AIRPORT
  - SELECT COUNT(Name) FROM NAME-ZIP
    WHERE City = 'Columbus'
    - Prints number of families in Columbus
    - Can also do other computations, like SUM, MIN, or AVG

- Result of a query is a subschema

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Security requirements

- Physical database integrity
- Logical database integrity
- Element integrity
- Referential integrity
- Auditability
- Access control
- User authentication
- Availability
Database integrity

- **Logical** and **physical** integrity
- Protect against database corruption
  - Allow only authorized individuals to perform updates
- Recover from physical problems (Power failures, disk crashes, . . .)
  - Perform periodic backups
  - Keep log of transactions to replay transactions since last backup

Element integrity

- Ensure correctness/accuracy of database elements
- **Access control** to limit who can update element
- **Element checks** to validate correctness
  - Element must be numeric, within a particular range, . . .
  - Not more than one employee can be president
  - Helps against mistakes by authorized users
  - Typically enforced by triggers (procedures that are automatically executed after an INSERT, DELETE, . . .)

Element integrity (cont.)

- **Change log** or **shadow fields** to undo erroneous changes
  - In case access control or element checks fail
  - Require additional space in the database
- **Error detection codes** to protect against OS or hard disk problems
Integrity: two-phase update

- For a set of operations, either all of them or none of them should be performed
  - Integrity violation if only some are performed
  - E.g., money is withdrawn from an account, but not deposited to another account
- First phase: gather information required for changes, but don’t perform any updates, repeat if problem arises
- Second phase: make changes permanent, repeat if problem arises
- See text for example

Integrity: concurrency control

- Concurrent modifications can lead to integrity violation
  - Two operations A and B read variable X
  - A then writes new value of X
  - B then writes new value of X
  - A’s update gets lost
- Need to perform A and B as atomic operations
- Take CS 454 for more about this

Referential integrity

- Each table has a primary key
- Minimal set of attributes that uniquely identifies each tuple
  - User ID or social insurance number
  - First name and last name (maybe not)
- A table might also have a or multiple foreign keys, which are primary keys in some other table
  - Zip is (likely) a primary key in ZIP-AIRPORT
  - Zip is a foreign key in NAME-ZIP
- Referential integrity ensures that there are no dangling foreign keys
  - For each zip in NAME-ZIP, there is an entry in ZIP-AIRPORT
Auditability

- Keep an audit log of all database accesses
  - Both read and write
- Access control can be difficult (see later), audit log allows to retroactively identify users who accessed forbidden data
  - Police officer looking at somebody’s criminal record as a favor to a friend, unauthorized medical personnel looking at Britney Spears’ medical records
- Maybe combination of accesses resulted in disclosure, not a single one (see later)
- Must decide about granularity of logging
  - Should results of a query be logged?

Access control

- More difficult than OS access control
- Might have to control access at the relation, record or even element level
- Many types of operations, not just read/write
  - SELECT, INSERT, UPDATE, CREATE, DROP,…
- Relationships between database objects make it possible to learn sensitive information without directly accessing it
  - Inference problem (see later)
- Efficiency problem in presence of thousands of records, each consisting of dozens of elements

Access control (cont.)

- Access control might consider past queries
  - Current query, together with past ones, could reveal sensitive information
    - Iteratively querying whether element is in set ultimately leaks set
- Or type of query
  - SELECT lastname, salary FROM staff
    WHERE salary > 50000
    might be forbidden, but not
  - SELECT lastname FROM staff
    WHERE salary > 50000
User authentication / Availability

- Database might do its own authentication
- Additional checks possible
  - E.g., time of day
- Databases facilitate sharing, but availability can suffer if multiple users want to access the same record
  - Block access until other user finishes updating record

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Types of data disclosure

- Exact data
- Bounds
  - Sensitive value is smaller than \( H \), but bigger than \( L \)
  - Might iteratively decrease range (binary search)
- Negative result
  - Knowing that a person does not have zero felony convictions is sensitive, even if actual number is hidden
- Existence
  - Knowing of existence of some data can be sensitive
- Probable value
  - Sensitive data has value \( x \) with probability \( y \)
Security vs. precision

• Security: Forbid any queries that access sensitive data, even if (aggregated) result is no longer sensitive
• Precision: Aggregated result should reveal as much non-sensitive data as possible

Data inference

• Derivation of sensitive data from (supposedly) non-sensitive data
• Direct attack
  • Attacker issues query that directly yields sensitive data
  • Might obfuscate query to fool DBMS
    • SELECT AVG(salary) FROM staff WHERE lastname = 'Adams'
    • OR (sex != 'M' AND sex != 'F')
• Indirect attack
  • Infer sensitive data from statistical results
    • As released by governments or pollers
  • Tracker attack

Statistical inference attacks

• Sum
  • Leaks sensitive data if sum covers only one record or if attacker can control set of covered records
    • SELECT SUM(salary)
    • SELECT SUM(salary) WHERE lastname != 'Adams'
• Count
  • Useful in attack above
• Mean
  • sum = count * mean
• Median
  • Intersecting medians might leak sensitive data
  • See text for example
Tracker attacks

• Focus on queries of type SUM or COUNT
• Assume that DBMS refuses to answer a query if
number of matching records is smaller than \( k \) or
larger than \( N - k \) and smaller than \( N \) (Why?)
  • \( N \): number of records in database
• A tracker \( T \) is a query whose result matches
between \( 2k \) and \( N - 2k \) records
  • DBMS will answer \( T \) (and not \( T \))
• Assume that there is a query \( C \) that DBMS
refuses to answer since it matches fewer than \( k \) or
more than \( N - k \) (but fewer than \( N \)) records

Tracker attacks (cont.)

• Let \( q() \) be the result of a query and \( S \) the set of
all records
• Using Venn diagrams, we can show that
  • \( q(C) = q(C \ or \ T) + q(C \ or \ not \ T) - q(S) \)
  • Use right-hand side for computing \( q(C) \) if \( q(C) \)
matches fewer than \( k \) records
  • \( q(C) = 2 \times q(S) - q(not \ C \ or \ T) - q(not \ C \ or \ not \ T) \)
  • Use right-hand side for computing \( q(C) \) if \( q(C) \)
matches more than \( n - k \) records
• In general, simple logic or linear algebra might
allow an attacker to convert a forbidden query into
multiple, allowed queries

Controls for statistical inference attacks

• Apply control to query or to data items
  • As seen, former is difficult
• Suppression and concealing are two controls
applied to data items
• Suppression
  • Suppress sensitive data from result
• Concealing
  • Answer is close to actual value, but not exactly
Controls (cont.)

- n-item k-percent rule
  - For the set of records that were included in the result, if there is a subset of n records that is responsible for over k percent of the result, omit the n records from result
  - However, omission itself might leak information or omitted value could be derived with other means

- Combined results
  - Report set or range of possible values

- Random sample
  - Compute result on random sample of database
  - Need to use same sample for equivalent queries

Controls (cont.)

- Random data perturbation
  - Add or subtract small random error to/from each value before computing result
  - Expectation is that statistical properties are maintained

- Query analysis
  - Maintain history of user’s queries and observe possible inferences
  - Costly, fails for colluding users

Data aggregation

- Data aggregation is related to data inference
- Building sensitive results from less sensitive inputs
- Aggregation can take place outside of a DBMS, which makes it difficult to control
  - People with different access rights talking to each other
- Closely related to data mining (see later), where information from different databases is combined
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Multilevel Security (MLS) Databases

- Support classification/compartmentalization of information according to its confidentiality/integrity
  - Two levels (sensitive and not sensitive) might not be sufficient
- At element level if necessary
  - Salary might be sensitive only for some employees
  - Other information in employee’s record might not be sensitive
- In an MLS database, each object has a classification and maybe a set of compartments
  - Object can be element, aggregate, column, or row

*-Property

- Implementing the *-property in an MLS database is difficult
  - User doing a write-up though user cannot read data at higher level (Blind writes)
  - Write-downs need a sanitization mechanism
  - Trusted processes that can do anything
- DBMS must have read and write access at all levels to answer user queries, perform back-ups, optimize database, . . .
  - Must trust DBMS
Confidentiality

- Depending on a user’s level, he/she might get different answers for a query
  - Less precision for low-level users
- Existence of a record itself could be confidential
- Keeping existence hidden can lead to having multiple records with the same primary key, but different sensitivity (polyinstantiation)
  - Admin notices that there is no record for employee Bob Hill and creates one
  - However, Bob Hill is a secret agent, so there already is a record, which admin cannot see
  - DBMS must allow admin’s request, else admin would get suspicious

Partitioning

- Have separate database for each classification level
- Simple, often used in practice
- Might lead to data stored redundantly in multiple databases
- Doesn’t address the problem of a high-level user needing access to low-level data combined with high-level data

Encryption

- Separate data by encrypting it with a key unique to its classification level
- Must be careful to use encryption scheme in the right way
  - E.g., encrypting the same value in different records with the same key should lead to different ciphertexts
- Processing of a query becomes expensive, many records might have to be decrypted
  - Doing the processing directly on the encrypted data is an active research area
Integrity lock

- Provides both integrity and access control
- Each data item consists of
  - The actual data item
  - A sensitivity label (maybe concealed)
  - A cryptographic signature (or MAC) covering the above plus the item’s attribute name and its record number
- Signature protects against attacks on the above fields, such as attacks trying to modify the sensitivity label, and attacks trying to move/copy the item in the database
- This scheme does not protect against replay attacks

Integrity lock (cont.)

- Any (untrusted) database can be used to store data items and their integrity locks
  - Locks can consume lots of space (maybe multiple locks per record)
- (Trusted) procedure handles access control and manages integrity locks
  - E.g., updates sensitivity level to enforce *-property or re-computes signature after a write access
  - Expensive
- Have to encrypt items and locks if there are other ways to get access to data in database
  - Makes query processing even more expensive

Trusted front end

- Front end authenticates a user and forwards user query to old-style DBMS
- Front end gets result from DBMS and removes data items that user is not allowed to see
- Allows use of existing DBMS and databases
- Inefficient if DBMS returns lots of items and most of them are being dropped by front end
Commutative filters

- Front end re-writes user query according to a user’s classification
  - Remove attributes that user is not allowed to see
  - Add constraint expressing user’s classification
- Benefits from DBMS’ superior query processing capabilities and discards forbidden data items early on
- Front end might still have to do some post processing

Distributed/federated databases

- Based on partitioning
- Front end forwards user query only to databases that user can access based on classification
- Front end might have to combine the results from multiple databases
  - Complex process, front end essentially becomes a DBMS
- Doesn’t scale to lots of classification levels

Views

- Many DBMS support views
- A view is logical database that represents a subset of some other database
  - CREATE VIEW foo AS SELECT * FROM bar WHERE...
- Element in view can correspond to an element in underlying database or be a combination of multiple elements
  - E.g., their sum
- Views can be used for access control
  - A user’s view of a database consists of only the data that the user is allowed to access
  - Hide attribute/row unless user is allowed to access at least one element, set to UNDEFINED any elements that user can’t access
Truman vs. non-Truman semantics

- **Truman semantics**: the DBMS pretends that the data the user is allowed to access is all the data there is
  - Like “The Truman Show”
  - All queries will succeed, even if they return imprecise results
- **Non-Truman semantics**: the DBMS can reject queries that ask for data the user is not allowed to access
  - Any queries that succeed will produce precise answers
  - Some queries will fail

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Data mining

- Multilevel databases weren’t a commercial success
  - Mainly military clients, finding all possible inferences is NP-complete
- However, the combination of (sensitive) information, stored in multiple (maybe huge) databases, as done for data mining, raises similar concerns and has gotten lots of attention recently
- So far, a single entity has been in control of some data
  - Knows what kind of data is available
  - Who has accessed it (ignoring side channels)
- No longer the case in data mining, data miners actively gather additional data from third parties
Data mining (cont.)

- Data mining tries to automatically find interesting patterns in data using a plethora of technologies
  - Statistics, machine learning, pattern recognition,…
  - Still need human to judge whether pattern makes sense (causality vs. coincidence)
- Data mining can be useful for security purposes
  - Learning information about an intrusion from logs

Security problems of data mining

- Confidentiality
  - Derivation of sensitive information
- Integrity
  - Mistakes in data
- Availability
  - (In)compatibility of different databases

Confidentiality

- Data mining can reveal sensitive information about humans (see later) and companies
- In 2000, the U.S. National Highway Traffic Safety Administration combined data about Ford vehicles with data about Firestone tires and become aware of a problem with the Ford Explorer and its Firestone tires
  - Problem started to occur in 1995, and each company individually had some evidence of the problem
  - However, data about product quality is sensitive, which makes sharing it with other companies difficult
- Supermarket can use loyalty cards to learn who buys what kind of products and sell this data, maybe to manufacturers’ competitors
Data correctness and integrity

- Data in a database might be wrong
  - E.g., input or translation errors
- Mistakes in data can lead to wrong conclusions by data miners, which can negatively impact individuals
  - From receiving irrelevant mail to being denied to fly
- Privacy calls for the right of individuals to correct mistakes in stored data about them
  - However, this is difficult if data is shared widely or if there is no formal procedure for making corrections
- In addition to false positives, there can also be false negatives: don’t blindly trust data mining applications

Availability

- Mined databases are often created by different organizations
  - Different primary keys, different attribute semantics, . . .
    - Is attribute “name” last name, first name, or both?
    - US or Canadian dollars?
- Makes combination of databases difficult
- Must distinguish between inability to combine data and inability to find correlation

Privacy and data mining

- Data mining might reveal sensitive information about individuals, based on the aggregation and inference techniques discussed earlier
- Avoiding these privacy violations is active research
- Data collection and mining is done by private companies
  - Privacy laws (e.g., Canada’s PIPEDA or U.S.’ HIPAA) control collection, use, and disclosure of this data
  - Together with PETs
- But also by governments
  - Programs tend to be secretive, no clear procedures
  - Phone tapping in U.S., no-fly lists in U.S. and Canada
Privacy-preserving data mining

- Anonymize data records before making them available
  - E.g., strip names, addresses, phone numbers
  - Unfortunately, such simple anonymization might not be sufficient
- August 6, 2006: AOL released 20 million search queries from 658,000 users
- To protect users’ anonymity, AOL assigned a random number to each user
  - 4417749 “numb fingers”
  - 4417749 “landscapers in Lilburn, Ga”
  - 17556639 “how to kill your wife”
- August 9: New York Times article re-identified user 4417749
  - Thelma Arnold, 62-year old widow from Lilburn, GA

Another example (by L. Sweeney)

- 87% of U.S. population can be uniquely identified based on person’s ZIP code, gender, and date of birth
- Massachusetts’ Group Insurance Commission released anonymized health records
- Records left away individuals’ name, but gave their ZIP code, gender, and date of birth (and health information, of course)
- Massachusetts’s voter registration lists contain these three items plus individuals’ names and are publicly available
- Enables re-identification by linking

\( k \)-anonymity [2002]

- Ensure that for each released record, there are at least \( k - 1 \) other released records from which record cannot be distinguished
- For health-records example, release a record only if there are \( k - 1 \) other records that have same ZIP code, gender, and date of birth
  - Assumption: there is only one record for each individual
- Because of the 87% number, this won’t return many records, need some pre-processing of records
  - Remove ZIP code, gender, or date of birth
  - Reduce granularity of ZIP code or date of birth
Discussion

- In health-records example, the attributes ZIP code, gender, and date of birth form a “quasi-identifier”
- Determining which attributes are part of the quasi-identifier can be difficult
  - Should health information be part of it?
  - Some diseases are rare and could be used for re-identification
  - However, including them is bad for precision
- Quasi-identifier should be chosen such that released records do not allow any re-identification based on any additional data that attacker might have
  - Clearly we don’t know all this data

Limitations of \( k \)-anonymity

A 3-anonymized table

<table>
<thead>
<tr>
<th>ZIP</th>
<th>DOB</th>
<th>Disease</th>
</tr>
</thead>
<tbody>
<tr>
<td>902**</td>
<td>196*-<em>-</em></td>
<td>Cancer</td>
</tr>
<tr>
<td>902**</td>
<td>196*-<em>-</em></td>
<td>Cancer</td>
</tr>
<tr>
<td>902**</td>
<td>196*-<em>-</em></td>
<td>Cancer</td>
</tr>
<tr>
<td>902**</td>
<td>195*-<em>-</em></td>
<td>Heart disease</td>
</tr>
<tr>
<td>902**</td>
<td>195*-<em>-</em></td>
<td>GI disease</td>
</tr>
<tr>
<td>902**</td>
<td>195*-<em>-</em></td>
<td>Flu</td>
</tr>
<tr>
<td>904**</td>
<td>195*-<em>-</em></td>
<td>Heart disease</td>
</tr>
<tr>
<td>904**</td>
<td>195*-<em>-</em></td>
<td>Cancer</td>
</tr>
<tr>
<td>904**</td>
<td>195*-<em>-</em></td>
<td>Cancer</td>
</tr>
</tbody>
</table>

\( \ell \)-diversity and \( t \)-closeness

- Homogeneity attack
  - If you know Bob (902**,196*-*-) is in the table, then Bob has cancer.
- Background knowledge attack
  - If you know Dave (904**,195*-*-) is in the table, and that his risk for heart disease is very low, then Dave has cancer.
- \( \ell \)-diversity property [2006]:
  - For any quasi-identifier, there should be at least \( \ell \) “well-represented” values of the sensitive fields
- Possibly still not good enough: \( t \)-closeness [2007]
  - Ensure that the distributions of the values for any quasi-identifier are within \( t \) of the distribution for the whole table

\( \Rightarrow \) Active research area
Value swapping

- Data perturbation based on swapping values of some (not all!) data fields for a subset of the records
  - E.g., swap addresses in subset of records
- Any linking done on the released records can no longer considered to be necessarily true
- Trade off between privacy and accuracy
- Statistically speaking, value swapping will make strong correlations less strong and weak correlations might go away entirely

Adding noise

- Data perturbation based on adding small positive or negative error to each value
- Given distribution of data after perturbation and the distribution of added errors, distribution of underlying data can be determined
  - But not its actual values
- Protects privacy without sacrificing accuracy

Recap

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