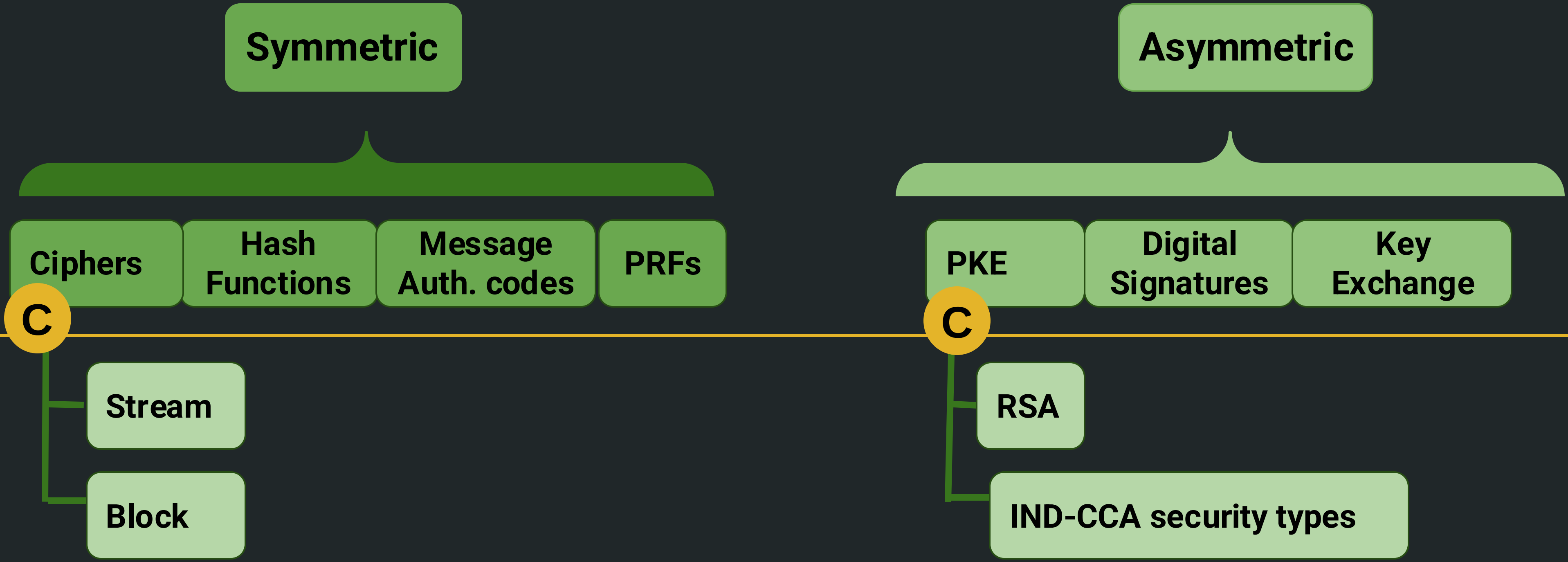
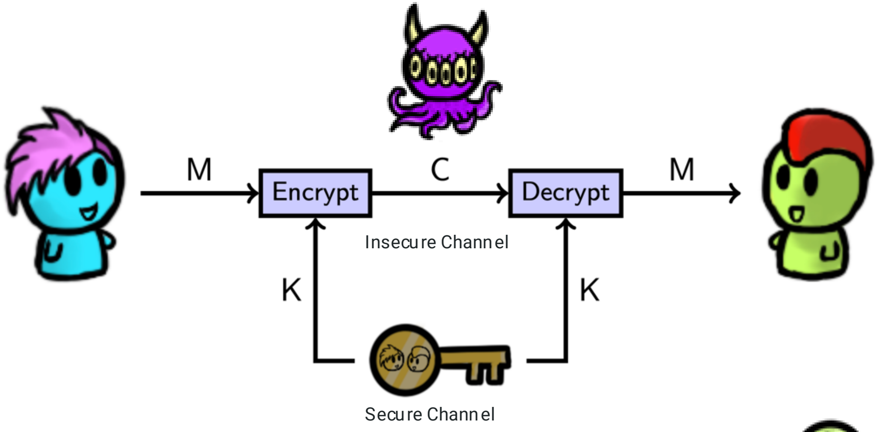
CS459/698 Privacy, Cryptography, Network and Data Security

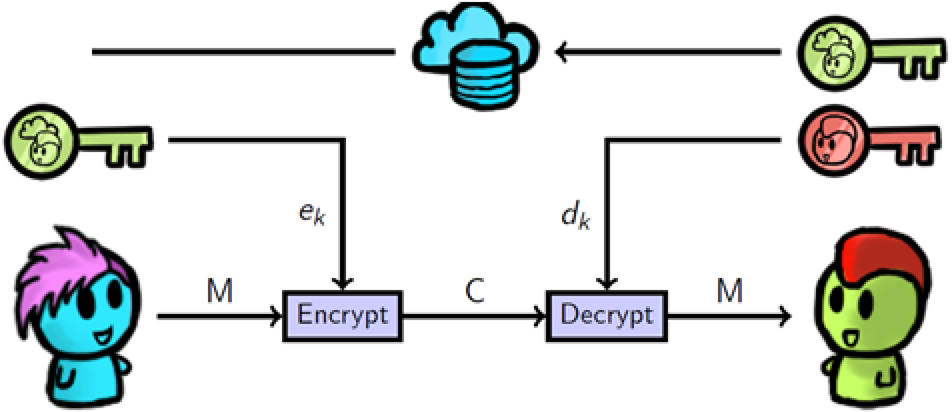
Integrity and Authenticated Encryption

Fall 2024, Tuesday/Thursday 02:30pm-03:50pm


Symmetric
Asymmetric
Ciphers
Hash Functions
Message Auth. codes
PRFs
PKE
Digital Signatures
Key Exchange
C
Stream
Block
C
RSA
IND-CCA security types

Block/Stream Ciphers, Public Key Cryptography…





Size of message on textbook RSA

- Overview:

$$(m^e)^d \equiv m \bmod N$$

Size of message on textbook RSA

- Overview:

$$(m^e)^d \equiv m \bmod N$$

m has to be strictly smaller than **N**, otherwise decryption will produce erroneous values.

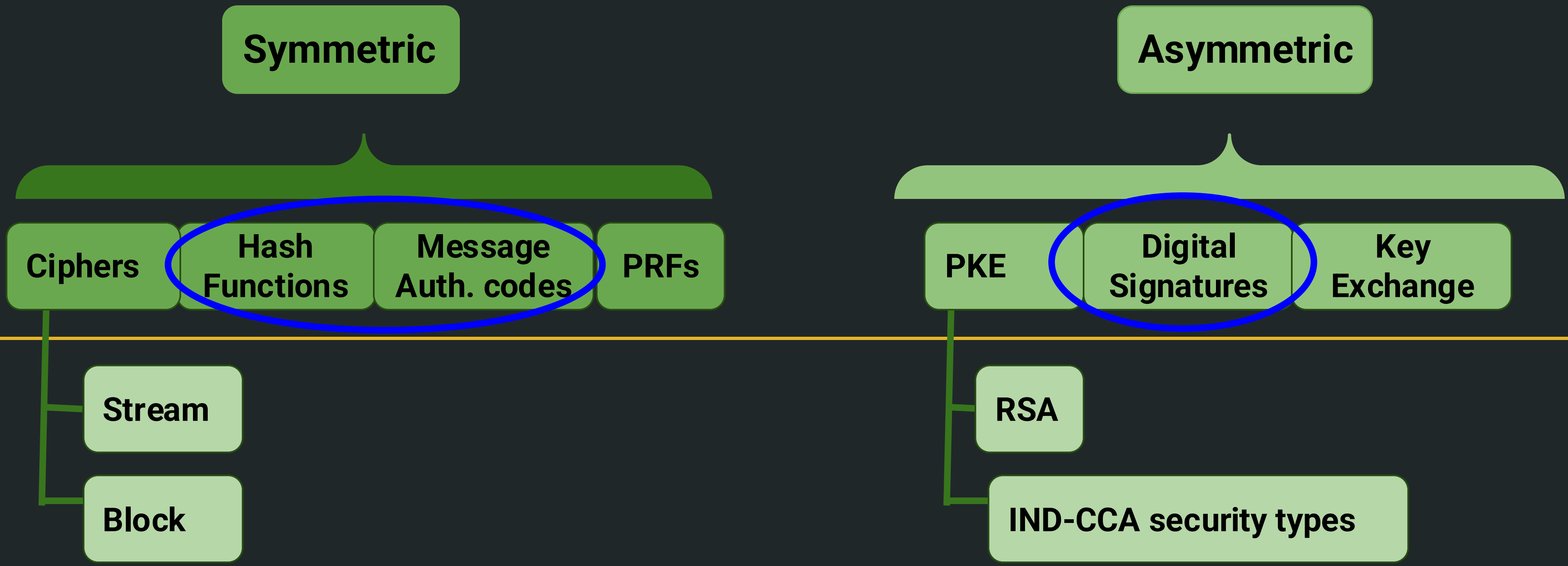
Size of message on textbook RSA

- Overview:

$$(m^e)^d \equiv m \bmod N$$

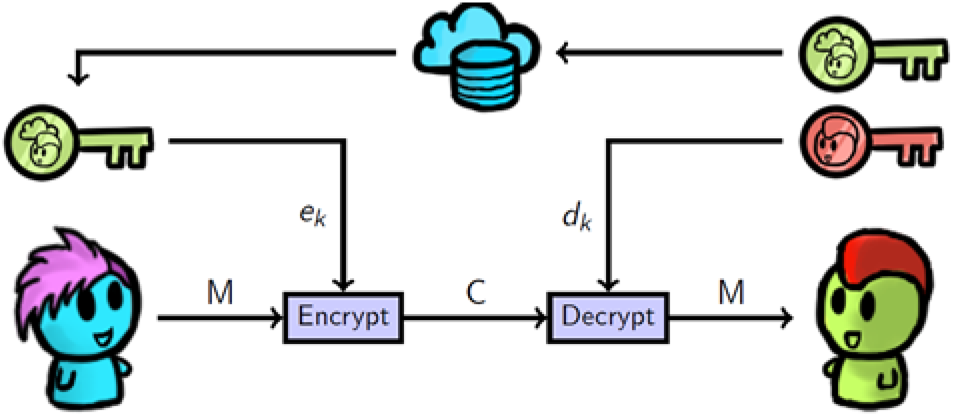
Is that all there is?

Goal: How do we make sure that Bob gets the same message Alice sent?


Symmetric
Asymmetric
Ciphers
Hash Functions
Message Auth. codes
PRFs
PKE
Digital Signatures
Key Exchange
Stream
Block
RSA
IND-CCA security types

Integrity components

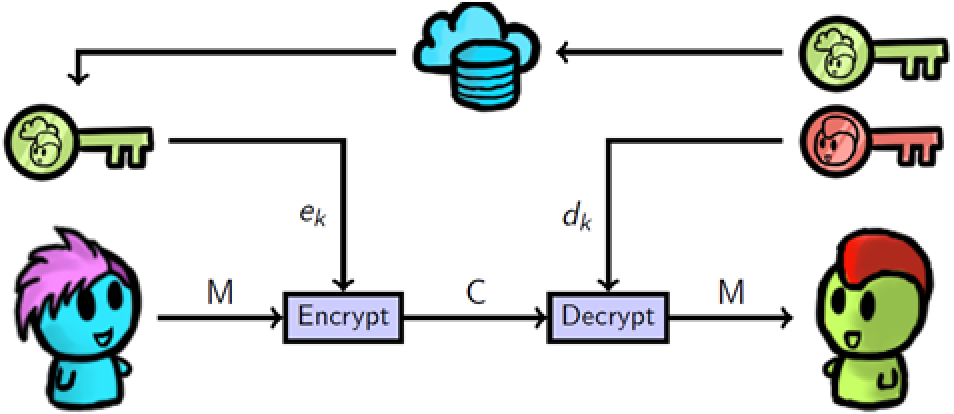
How do we tell if a message has changed in transit?



…wait…is this the message Alice sent?

Integrity components

How do we tell if a message has changed in transit?



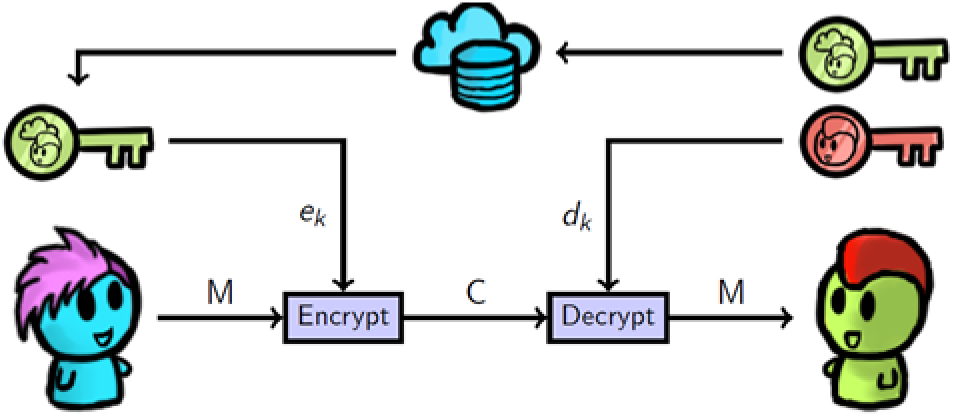
…wait…is this the message Alice sent?

Checksums

ISBN

Integrity components

How do we tell if a message has changed in transit?



…wait…is this the message Alice sent?

Checksums

Add up all the bytes of M, append the checksum to M so Bob can verify it

Not. Good. Enough.

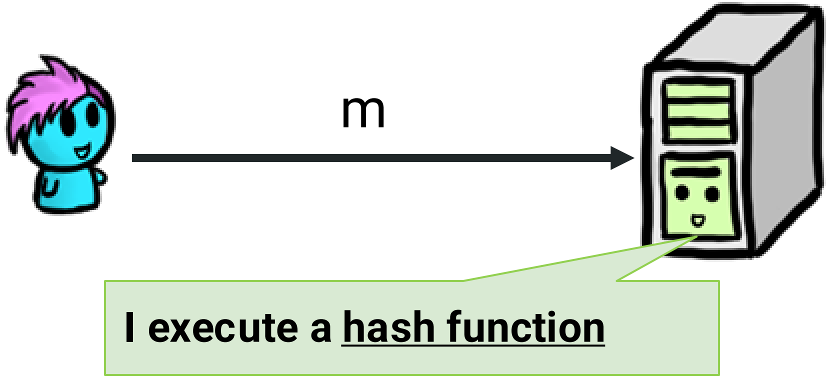
Checksums are deterministic…

Not. Good. Enough.

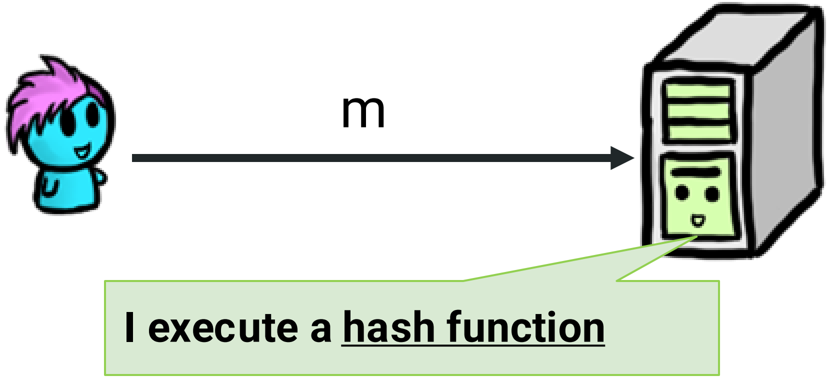
Goal: Make it hard for Mallory to find a second message with the same checksum as the “real” message

“Cryptographic” checksum

Cryptographic hash functions

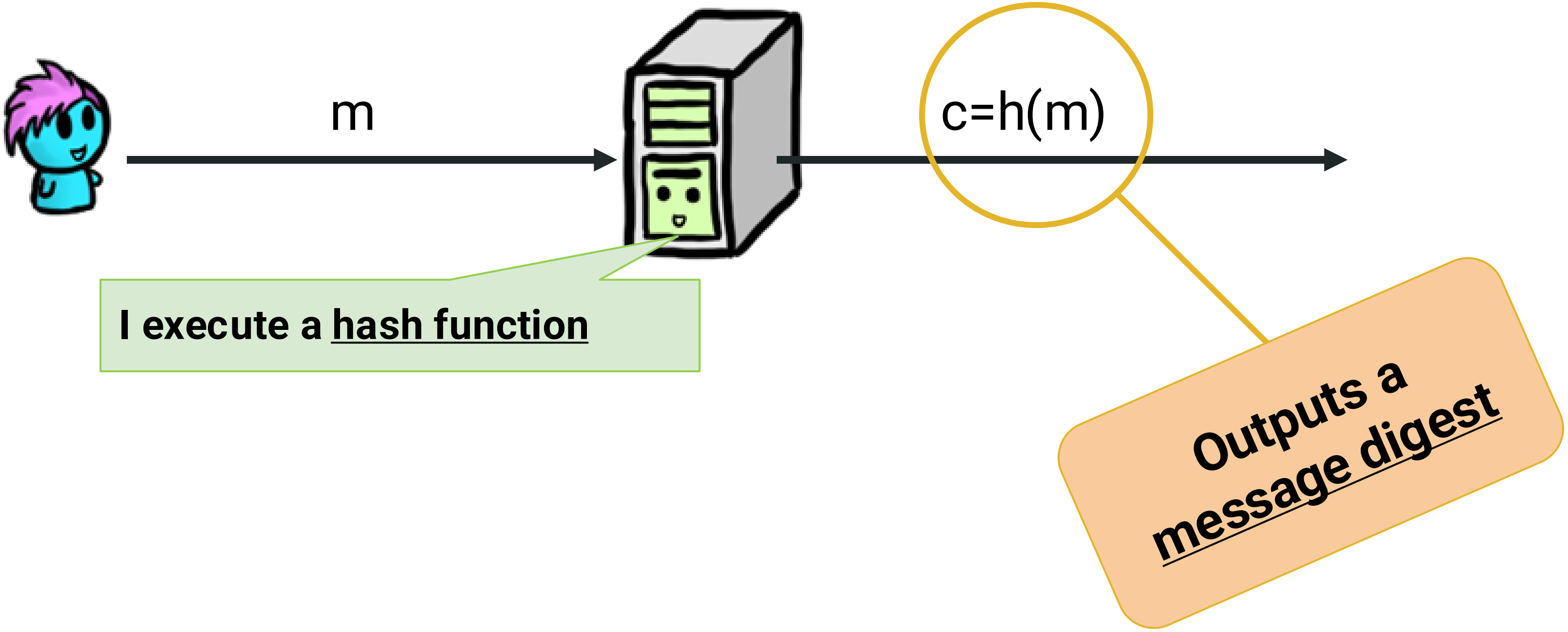


Cryptographic hash functions

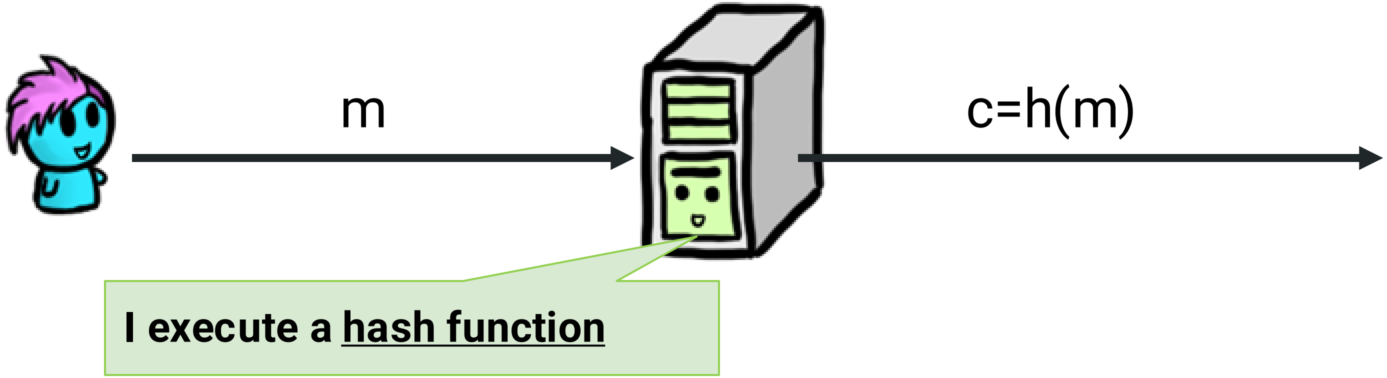


Takes an arbitrary length string, and computes a fixed length string.

Cryptographic hash functions



Cryptographic hash functions

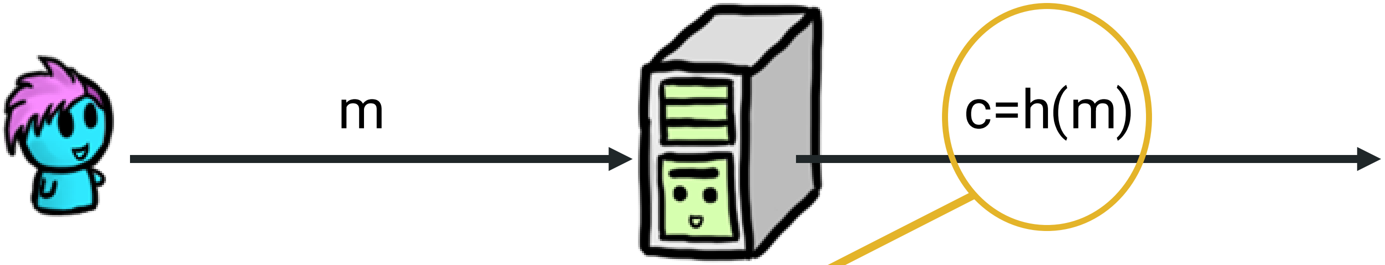


Q: Why is this useful?

Common examples:

- MD5, SHA-1, SHA-2, SHA-3 (aka Keccak after 2012)

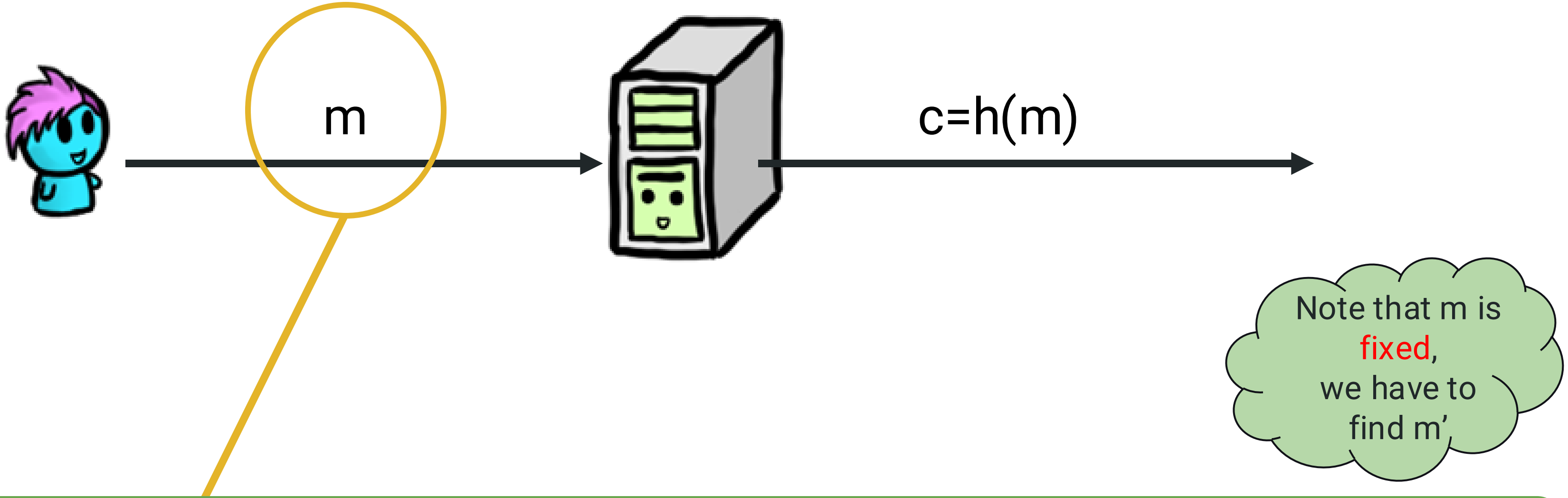
Properties: Preimage-Resistance



Goal: Given c, it's "hard" to find m such that h(m) = c

(i.e., a "preimage" of x)

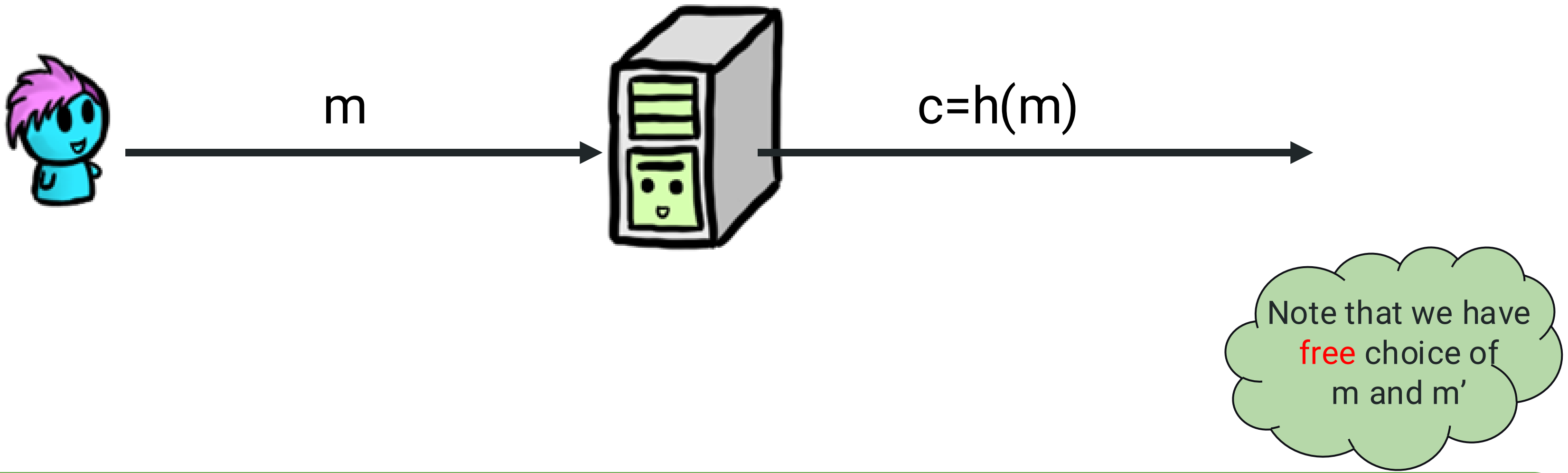
Properties: Second Preimage-Resistance



Goal: Given m, it's "hard" to find m' ≠ m such that h(m) = h(m')

(i.e., a "second preimage" of h(x))

Properties: Collision-Resistance



Goal: It’s hard to find any two distinct m, m’ such that h(m) = h(m’)
(i.e., a “collision”)

What do we mean by “hard”?

- SHA-1: takes 2^{160} work to find a preimage or second image

- SHA-1: takes 2^{80} to find a collision using brute-force search
  - However, there are faster ways than brute-force to find collisions in SHA-1 or MD5
  - → For a hash function with an n-bit output, the birthday attack can find collisions in approximately $2^{n/2}$ computations. (2^{80} evaluations)

MD5 2008 collision attack against digital certificates.

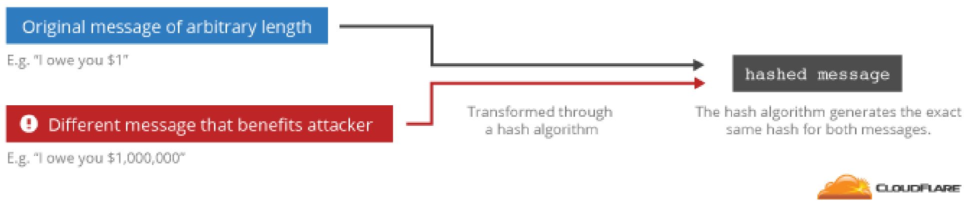
Create 2 cert. with same MD5 hash

SHA-1 2017 collision attack against digital certificates.

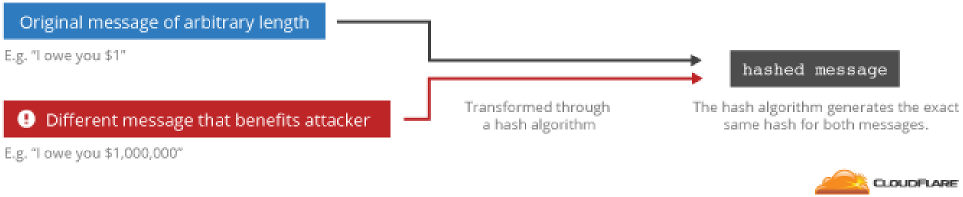
Making it too hard to break these properties?

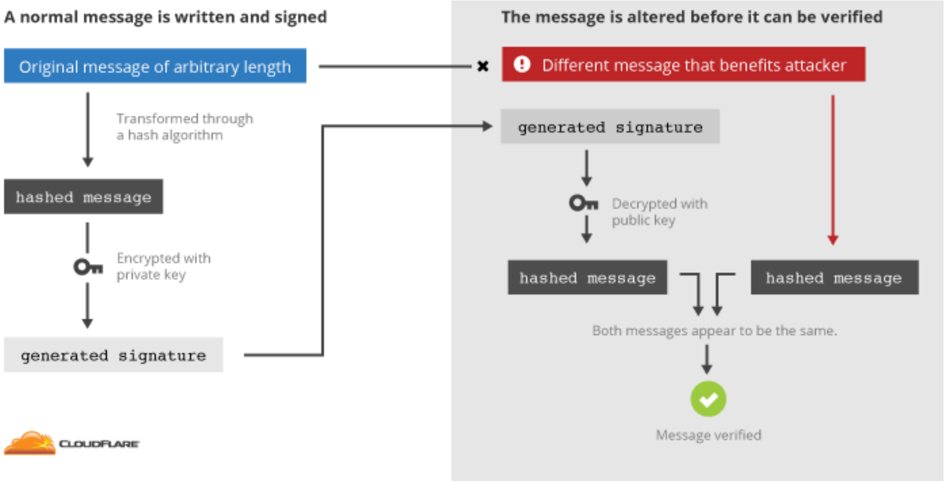
- SHA-1: takes 2^{160} work to find a preimage or second image
- SHA-1: takes 2^{80} to find a collision using brute-force search
    - However, there are faster ways than brute-force to find collisions in SHA-1 or MD5
- Collisions are always easier to find than preimages or second preimages due to the birthday paradox

How collisions work



How attackers exploit hash collisions





The birthday paradox

- If there are **n** people in a room, what is the probability that at least two people have the same birthday?
- For n = 2: Pr(2) = 1 – $\frac{364}{365}$
- For n = 3: Pr(3) = 1 – $\frac{364}{365}$ x $\frac{363}{365}$
- For n people: Pr(n) = 1 – $\frac{364}{365}$ x $\frac{363}{365}$ x … x $\frac{365-n-1}{365}$

Collisions and the Birthday Paradox

Collisions are easier due to the birthday paradox

What’s the probability two of us have the same birthday?

Collisions and the Birthday Paradox

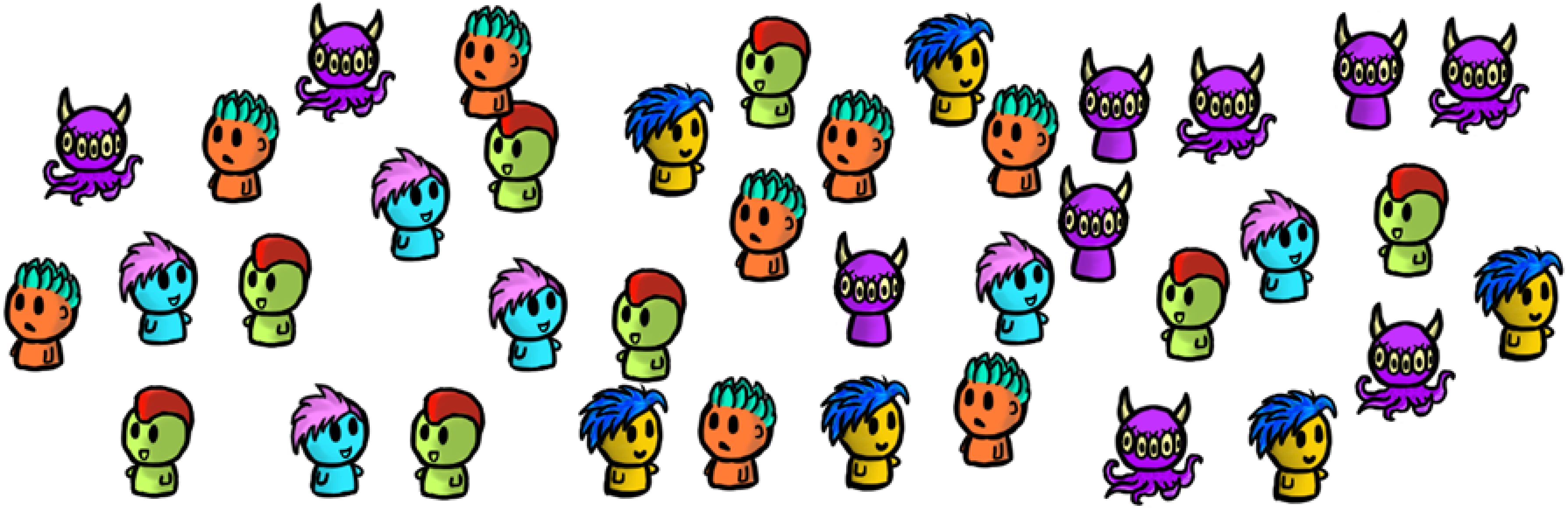
Collisions are easier due to the birthday paradox

What's the probability two of us have the same birthday?

There's 23 of us, so larger than 50%!!

Collisions and the Birthday Paradox

Collisions are easier due to the birthday paradox



Collisions and the Birthday Paradox

Collisions are easier due to the birthday paradox

Collisions and the Birthday Paradox

Collisions are easier due to the birthday paradox

Collisions and the Birthday Paradox

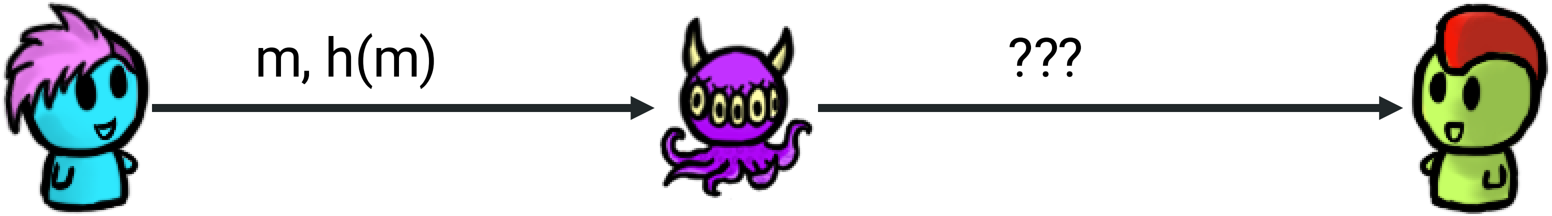
Collisions are eas… due t…

Th… 9%!!!

This is NOT the end of our problems…

How about a bad example?

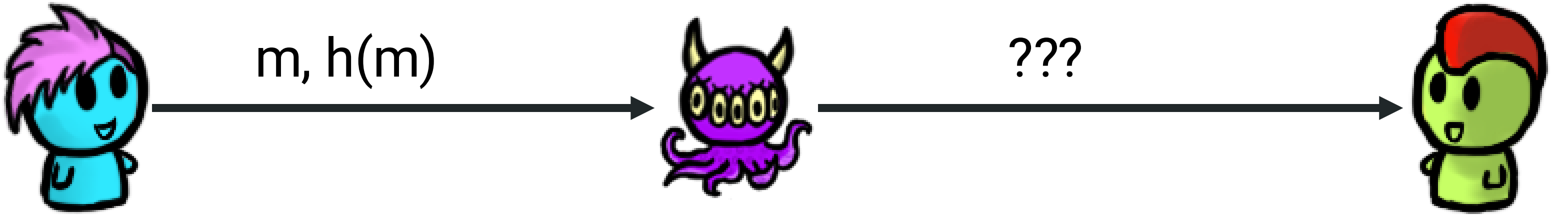
Assume we don't care about confidentiality now, just integrity.



Q: What can Mallory do to send the message she wants (change m)?

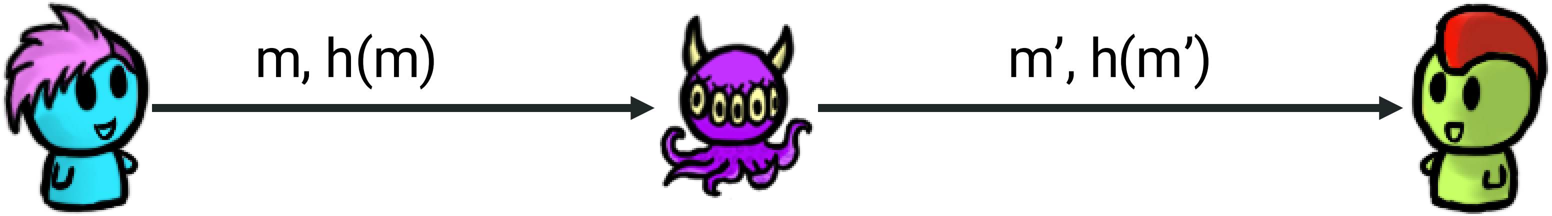
How about a bad example?

Assume we don’t care about confidentiality now, just integrity.



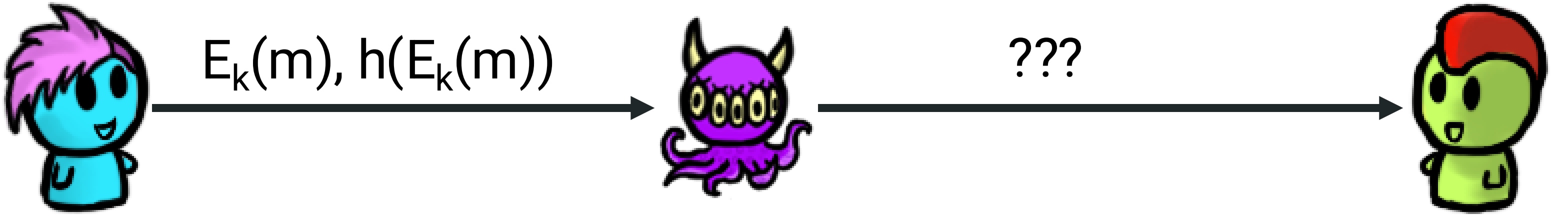
Q: What can Mallory do to send the message she wants (change m)?

A: Just change it…Mallory can compute the new hash herself.



How about a bad example?

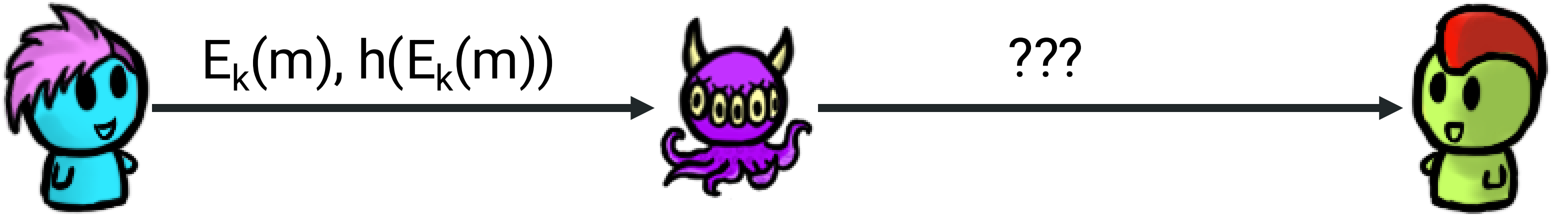
Assume we also care about confidentiality now.



Q: What can Mallory do to send the message she wants (change E(m))?

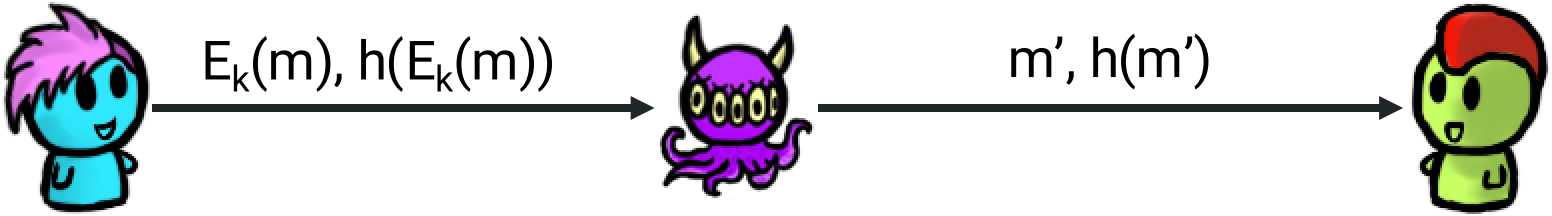
How about a bad example?

Assume we also care about confidentiality now.



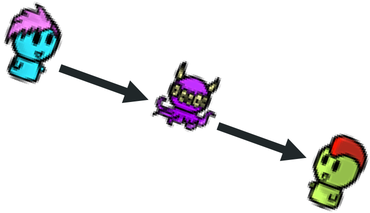
Q: What can Mallory do to send the message she wants (change E(m))?

A: Still just change it.



Limitations for Cryptographic Hash Functions

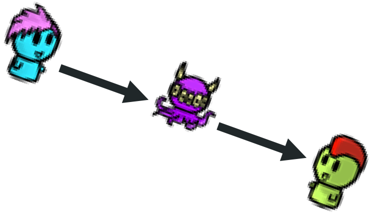
- Integrity guarantees only when there is a **secure** way of sending/storing the message digest



I could publish the hash of my public key on a business card

Limitations for Cryptographic Hash Functions

- Integrity guarantees only when there is a **secure** way of sending/storing the message digest

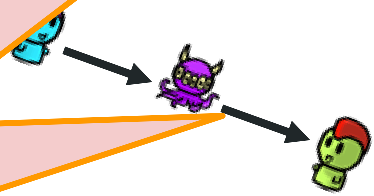


I could publish the hash of my public key on a business card

Good idea! Although the key would be too big to place on the card, I could use the hash to… verify it!

Limitations for Cryptographi Hash Functions

- Integrity guarantees only whe
way of sending/sto the



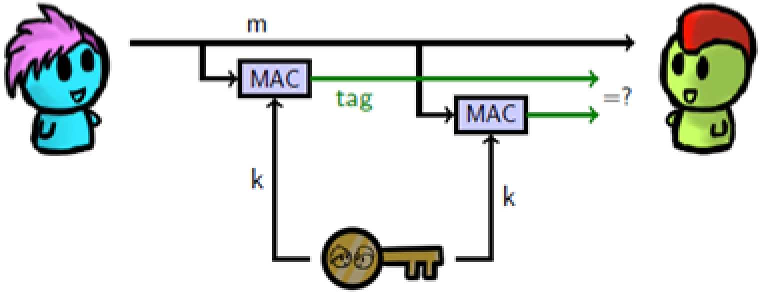
I could publish the hash of my public key o business card

the key would
card, I
e the hash to... verify it!

What if…we don't have an external/physical channel? i.e., using the Internet to communicate

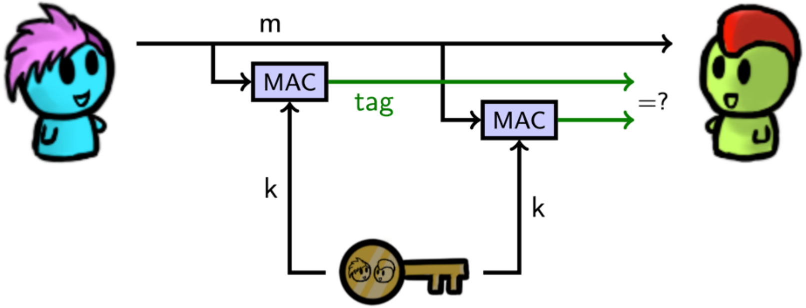
Authentication and Hash Functions

- We can use “keyed hash functions”
- Requires a secrete key to generate, or even check, the computed hash value (sometimes called a tag)



Called: Message authentication codes (MACs)

Message Authentication Codes (MACs)



I don't have the key to generate or check the values…

Do the MAC/tag values match?

YES	NO
No one messed with the data	The data has been altered somehow

Common examples:

- SHA-1-HMAC, SHA-256-HMAC, CBC-MAC

Combine Ciphers and MACs

Confidentiality

Integrity

Combine Ciphers and MACs

Confidentiality

+

Integrity

In practical we often need both confidentiality and message integrity

But how to combine them? Three possibilities

There are multiple strategies to combine a cipher and a MAC when processing a message

MAC-then-Encrypt, Encrypt-and-MAC, Encrypt-then-MAC

But how to combine them? Three possibilities

There are multiple strategies to combine a cipher and a MAC when processing a message

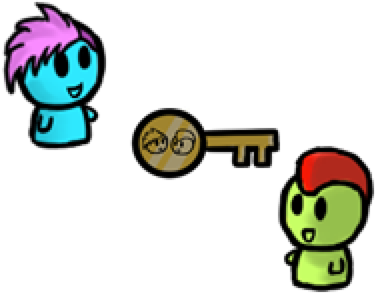
MAC-then-Encrypt, Encrypt-and-MAC, Encrypt-then-MAC

Ideally crypto libraries already provides an authenticated encryption mode that securely combines the two operations, so we don't have to worry about getting it right

- E.g., GCM, CCM (used in WPA2, see later), or OCB mode

Let’s try it!

- Alice and Bob have a secret key **K** for a cryptosystem ($E_k(\cdot)$, $D_k(\cdot)$)

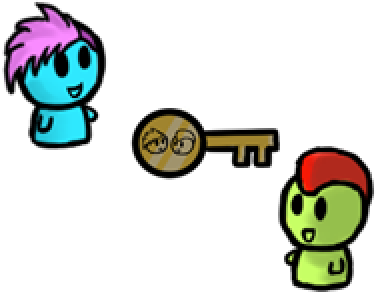


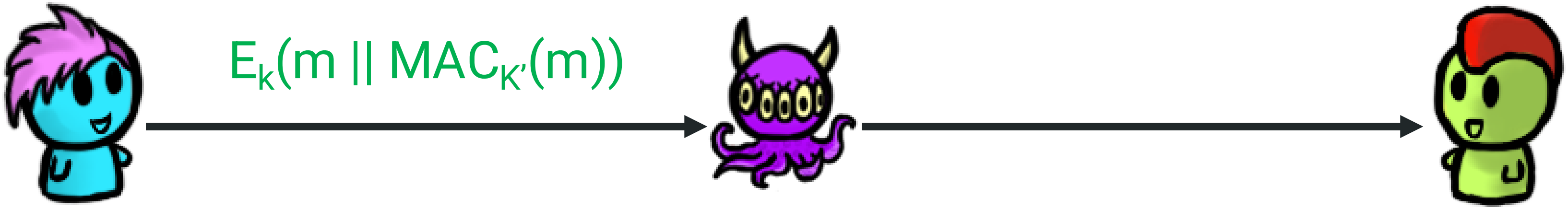
- Also, a secret key **K’** for their $MAC_{K'}(\cdot)$

How can Alice build a message for Bob in the following three scenarios?

MAC-then-Encrypt

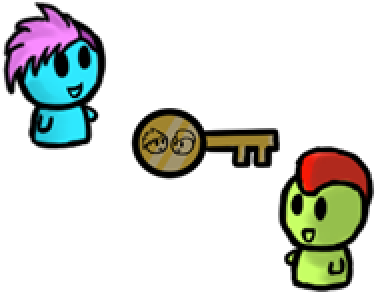
- Compute the MAC on the message, then encrypt the message and MAC together, and send that ciphertext.

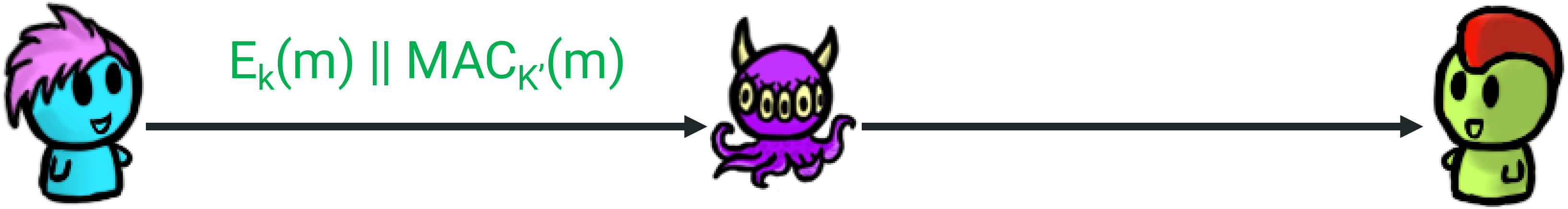




Encrypt-and-MAC

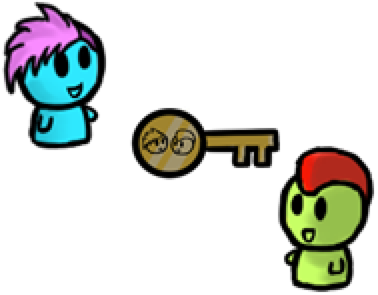
- Compute the MAC on the message, the encryption of the message, and send both.

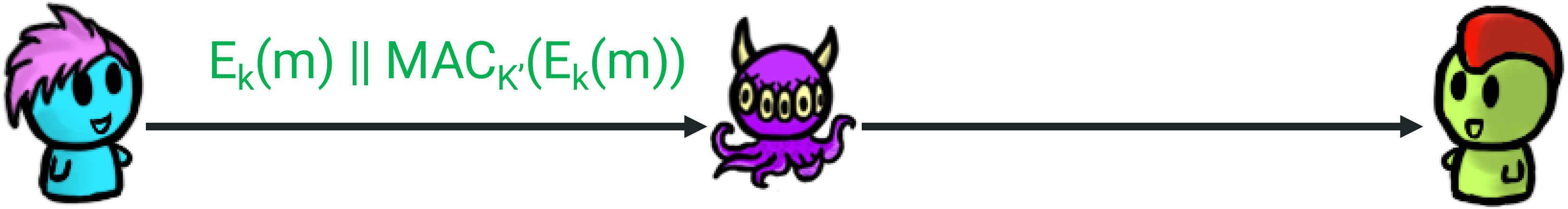




Encrypt-then-MAC

- Encrypt the message, compute the MAC on the encryption, send encrypted message and MAC





Which order is correct?

Q: Which should be recommended then?

$E_k(m \,||\, MAC_{K'}(m))$ **vs.** $E_k(m) \,||\, MAC_{K'}(m)$ **vs.** $E_k(m) \,||\, MAC_{K'}(E_k(m))$

MAC-then-encrypt | **Encrypt-and-MAC** | **Encrypt-then-MAC**

The Doom Principle

"if you have to perform any cryptographic operation before verifying the MAC on a message you've received, it will somehow inevitably lead to doom."

The Doom Principle

"if you have to perform any cryptographic operation before verifying the MAC on a message you've received, it will somehow inevitably lead to doom."

Q: What are possible problems that can arise from the orderings?

The Doom Principle

Q: What are possible problems that can arise from the orderings?

- **MAC-then-Encrypt:** Allows an adversary to force Bob into decrypting the ciphertext before verifying the MAC. May lead to a padding oracle attack

The Doom of MAC-then-Encrypt

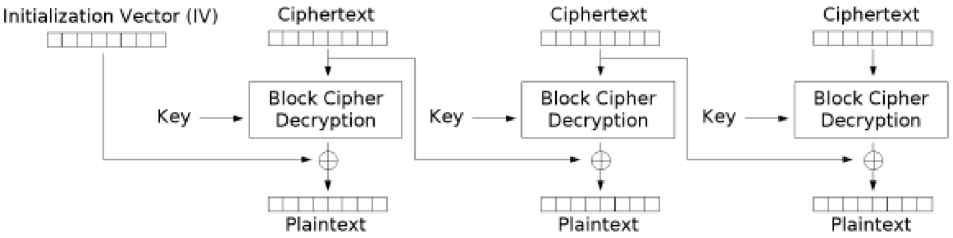
Observation: To verify the MAC, Bob has first to decrypt the message, since the MAC is part of the encrypted payload

- **Padding oracle attack:** The idea is for the attacker to send modified ciphertexts to Bob and observe how he responds.
- With CBC, by modifying the last block of the ciphertext in a way that alters the block's padding, the attacker can tell if the padding is valid or not.
- If the padding is invalid, the system might respond differently (e.g., with an error message that is padding-specific). This information leakage allows the attacker to gradually decrypt the ciphertext byte by byte.

The Doom of MAC-then-Encrypt

Observation: To verify the MAC, Bob has first to decrypt the message, since the MAC is part of the encrypted payload

- **Padding oracle attack:**
  - So if a block needs to be padded out by 5 bytes, for instance, one would append 5 bytes of the value 0x05.
  - 1st decrypt the message, look at the value of the last byte (call it N), and then insure that the preceding N-1 bytes also had the value of N.
  - If we encounter an incorrect value → padding error, and should abort. Since the MAC is part of the encrypted payload, *all* of this needs to happen before the MAC can be verified.



Cipher Block Chaining (CBC) mode decryption

The Doom Principle

Q: What are possible problems that can arise from the orderings?

- **Encrypt-and-MAC:** Allows an adversary to force Bob into decrypting the ciphertext to verify the MAC. May lead to a chosen-ciphertext attack

The Doom of Encrypt-and-MAC

Q: What happens if the MAC has no mechanism to provide confidentiality?

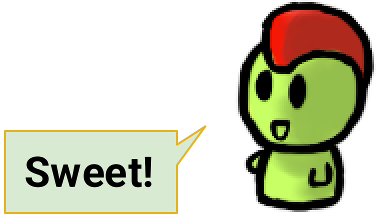
- MACs are meant to provide integrity
- MACs are often implemented by a **deterministic** algorithm without an explicit random input (essentially, for a given key and message, the output of the MAC is always the same).
- If a deterministic MAC is used, then there is no guarantee that the tag $E_k(m) \| \mathbf{MAC_{K'}(m)}$ will not leak information about the secret message **m**.

Which order is correct?

We want the receiver to verify the MAC first!

The recommended strategy is Encrypt-then-MAC:
$E_k(m) \,||\, MAC_{K'}(E_K(m))$

- **Encrypt-then-MAC:** Allows Bob to check the MAC of the ciphertext before performing any decryption whatsoever (e.g., prevent attacks by immediately closing a connection if the MAC fails)



More properties that matter?

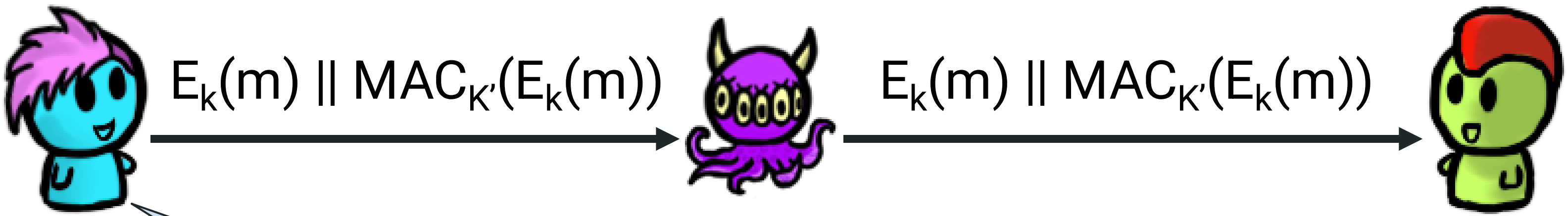
Repudiation

Alice sent *m*, and I received the same m she sent.

Repudiation

Confidentiality **Integrity** **Authentication**

Repudiation



Almost, but not quite a "signature"

Confidentiality **Integrity** **Authentication**

Repudiation

So…you’re saying Bob can’t prove to Carol that Alice sent m?

Repudiation

So…you're saying Bob can't prove to Carol that Alice sent m?

Q: Why can't Bob prove it?

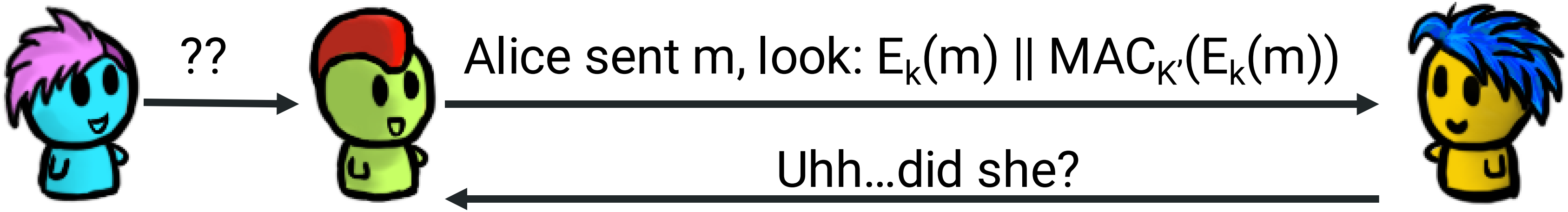
Repudiation

So…you're saying Bob can't prove to Carol that Alice sent m?

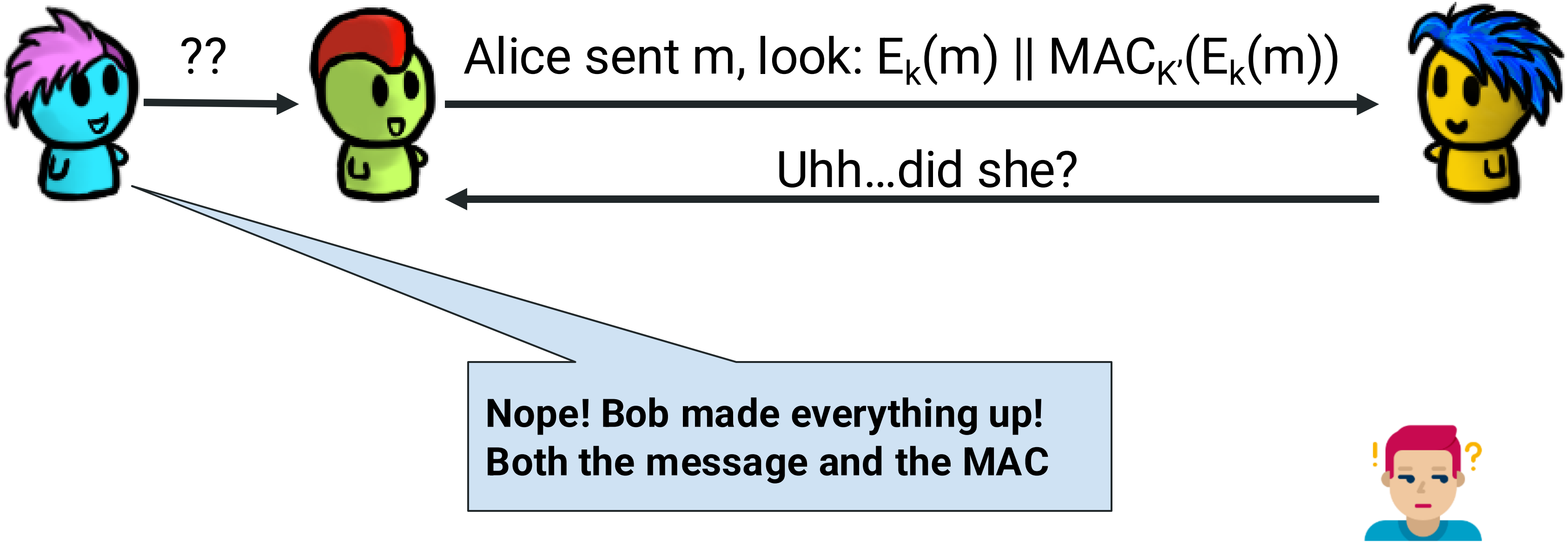
Q: Why can't Bob prove it?

A: Either Alice or Bob could create any message and MAC combination…also Carol doesn't know the secret keys.

Implications?

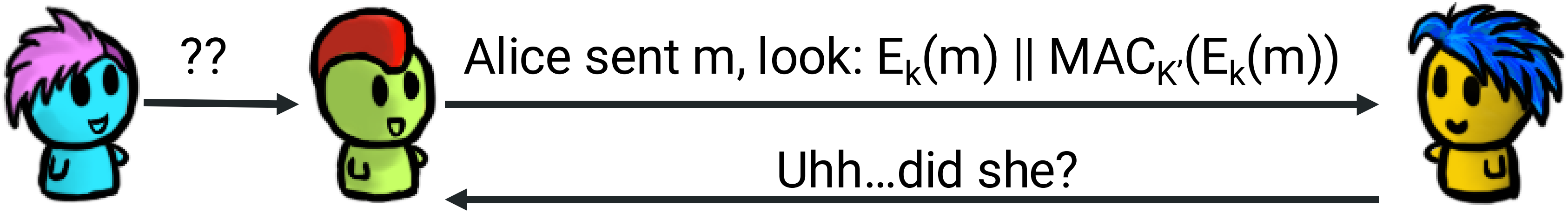


Implications?



Bob be like

Implications?

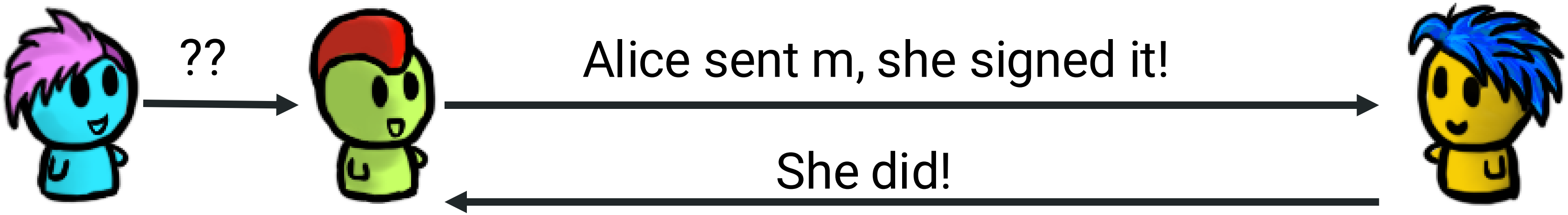


This is called repudiation, and we sometimes want to avoid it

Repudiation Property: For some applications this property is good (e.g., private conversations)…others less good (e.g., e-commerce…).

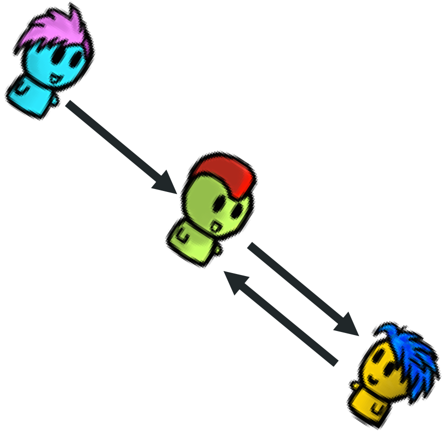
Digital Signatures

For non-repudiation, what we want is a true digital signature, with the following properties:



Properties of digital signatures

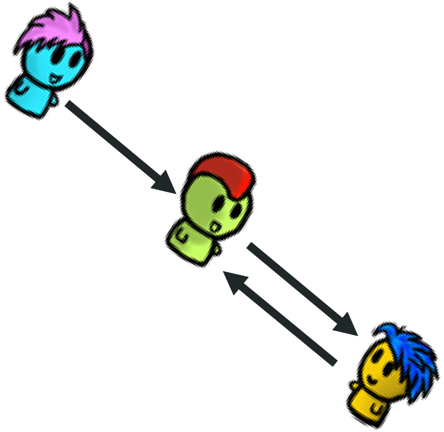
If Bob receives a message with Alice's digital signature on it, then:



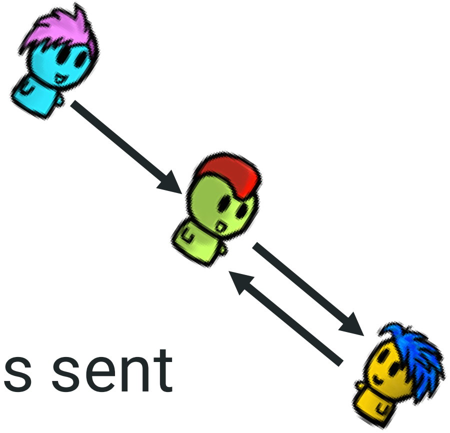
Properties of digital signatures

If Bob receives a message with Alice's digital signature on it, then:

- Alice sent it, and not , (like a MAC)



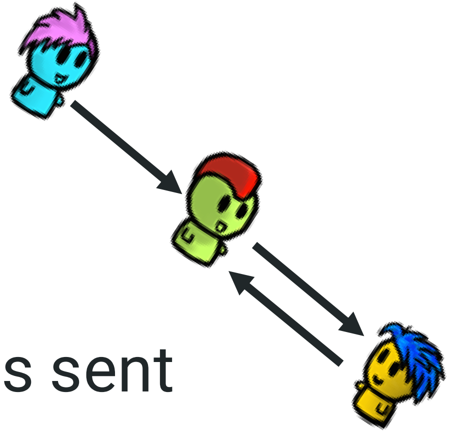
Properties of digital signatures



If Bob receives a message with Alice's digital signature on it, then:

- Alice sent it, and not , (like a MAC)
- The message has not been altered since it was sent (like a MAC)

Properties of digital signatures

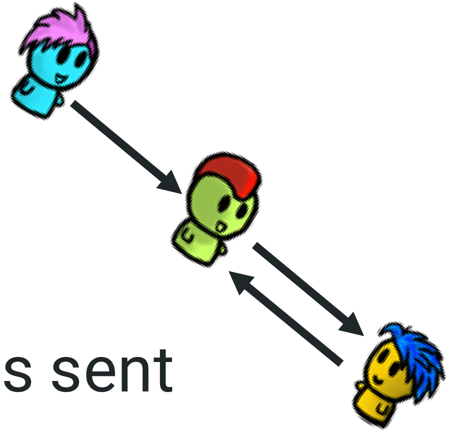


If Bob receives a message with Alice’s digital signature on it, then:

- Alice sent it, and not , (like a MAC)
- The message has not been altered since it was sent (like a MAC)
- Bob can prove these properties to a third party (NOT like a MAC)

Properties of digital signatures

If Bob receives a message with Alice’s digital signature on it, then:



- Alice sent it, and not , (like a MAC)
- The message has not been altered since it was sent (like a MAC)
- Bob can prove these properties to a third party (NOT like a MAC)

Achievable? Use techniques similar to public-key crypto (last class)

Making Digital Signatures

1. A pair of keys 

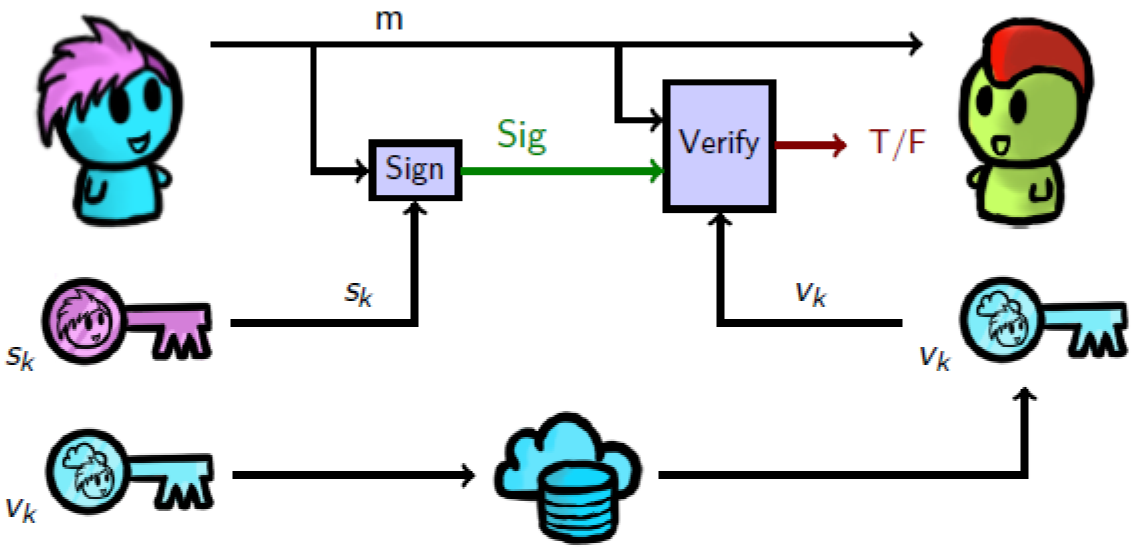
2. Everyone gets Alice public verification key 

3. Alice signs m with her private signature key S_k 

4. Bob verifies m with Alice's public verification key V_k 

5. If it verifies correctly, the signature is valid

Digital Signatures at a Glance



Faster Signatures

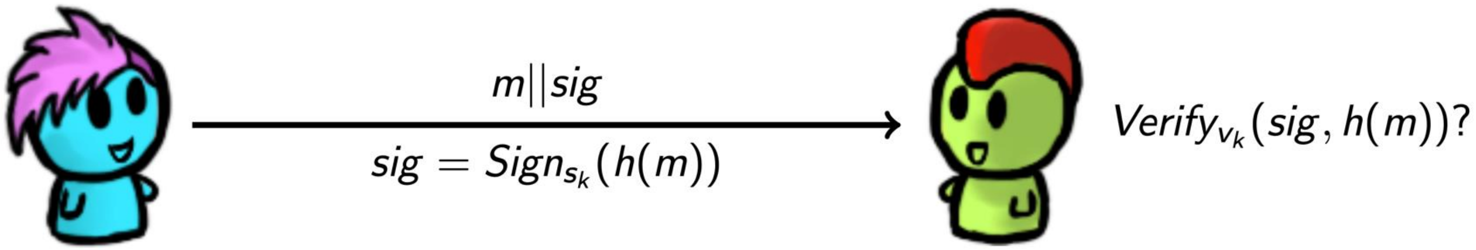
- Signing large messages is slow

  → "hybridize" the signatures to make them faster
- A hash is much smaller than the message… faster to sign

Faster Signatures

- Signing large messages is slow

  → "hybridize" the signatures to make them faster
- A hash is much smaller than the message… faster to sign

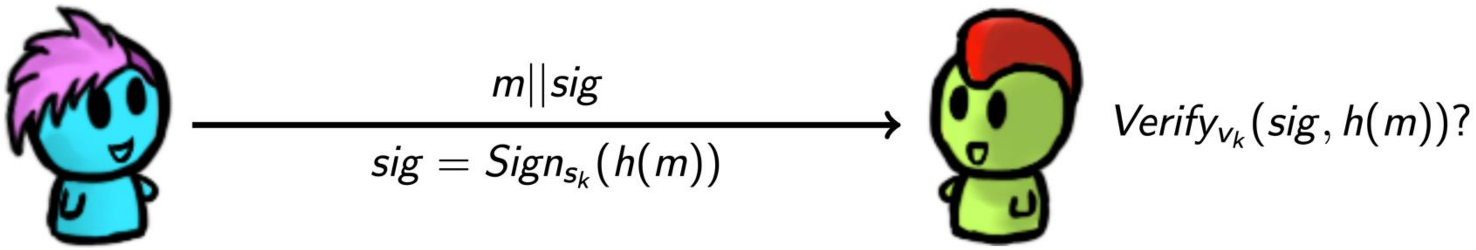


Faster Signatures

- Signing large messages is slow

  → "hybridize" the signatures to make them faster

- A hash is much smaller than the message… faster to sign



- Finally, authenticity and confidentiality are separate

  → you need to include both if you want to achieve both

Combining PKE and digital signatures

- Alice has two different key pairs:
  - → an (encryption, decryption) key pair e_k^A, d_k^A
  - → an (signature, verification) key pair s_k^A, v_k^A
- So does Bob : e_k^B, d_k^B and s_k^B, v_k^B
- Alice uses e_k^B to encrypt a message destined for Bob:
  - → $C = E_{e_k^B}(M)$
- She uses s_k^A to sign the ciphertext:
  - → $T = Sign_{s_k^A}(C)$
- Bob uses v_k^A to check the signature:
  - → $Verify_{v_k^A}(C,T)$, if verified, C is authentic
- He uses d_k^B to check the signature:
  - → $M = D_{d_k^B}(C)$

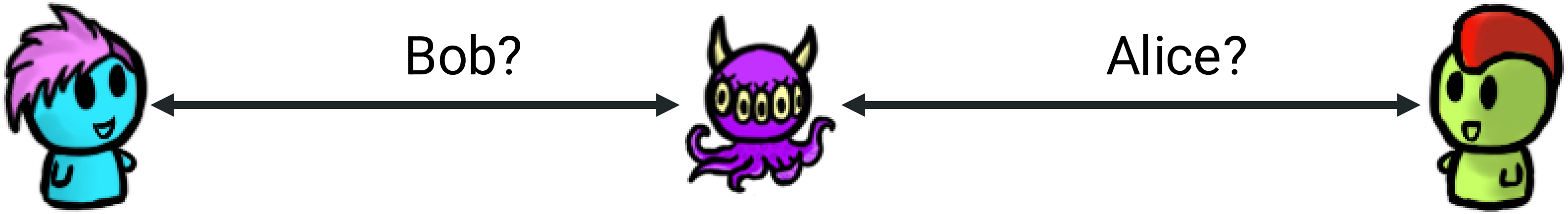
Relationship between key pairs

- Alice (signature, verification) key pair is long-lived, whereas her (encryption, decryption) key pair is short-lived

  → Provides forward secrecy

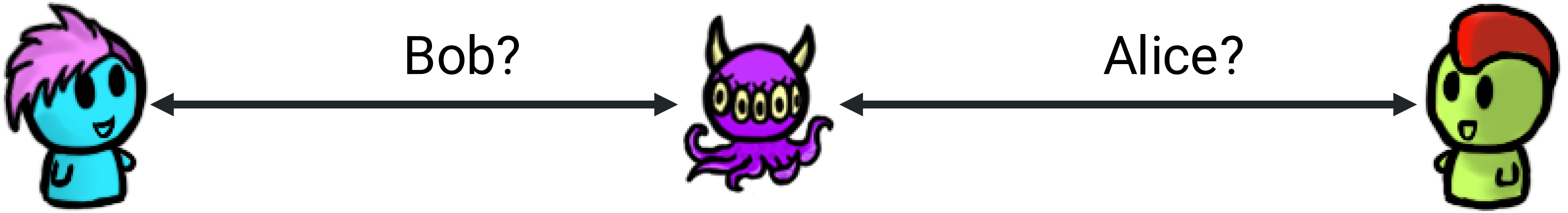
- When creating a new (encryption, decryption) key pair, Alice uses her signing key to sign her new encryption key and Bob uses Alice's verification key to verify the signature on this new key

The Key Management Problem



Q: How can Alice and Bob be sure they're talking to each other?

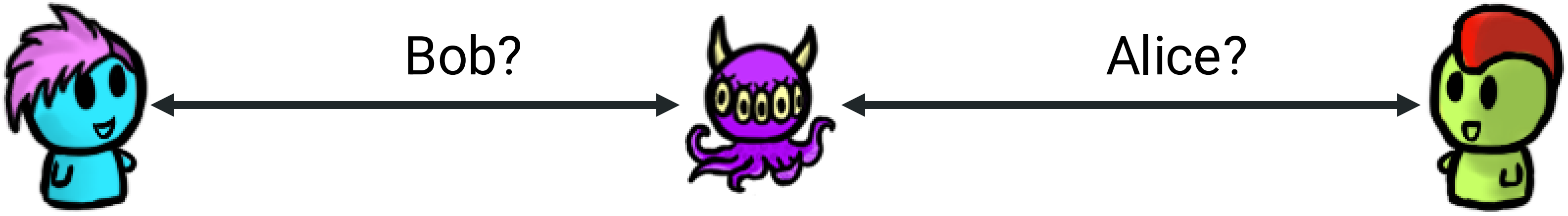
The Key Management Problem



Q: How can Alice and Bob be sure they’re talking to each other?

A: By having each other’s verification key!

The Key Management Problem

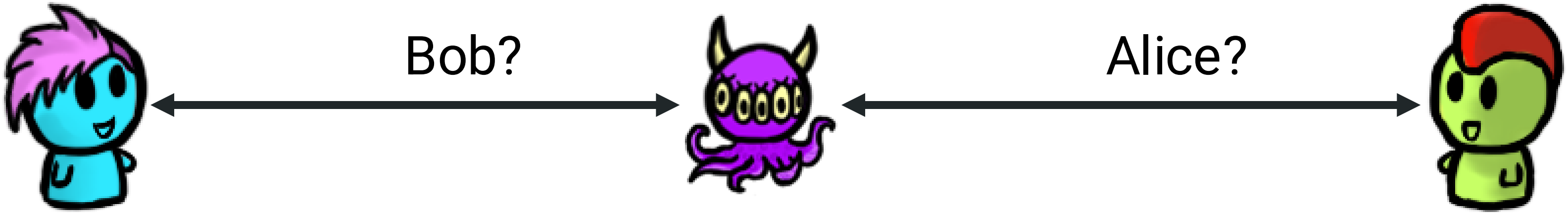


Q: How can Alice and Bob be sure they're talking to each other?

A: By having each other's verification key!

Q: But how do they get the keys…

The Key Management Problem…Solutions?



Q: But how do they get the keys…

A: Know it personally (manual keying e.g., SSH)

A: Trust a friend (web of trust e.g, PGP)

A: Trust some third party to tell them (CAs, e.g., TLS/SSL)

Nex up: More Cryptography…